

Semester I
Winter Examinations 2007/2008

Exam Code(s)	3IF1
Exam(s)	IF1 B.Sc. (Information Technology)
Module Code(s)	CT306
Module(s)	Formal Methods
Paper No.	1
Repeat Paper	Special Paper
External Examiner(s)	Professor J. Keane
Internal Examiner(s)	Professor. G. Lyons
	Dr. M. Mc Gettrick

Instructions

Answer 3 questions.
All questions will be marked equally.

Duration	2hrs
No. of Answer Books	1

Requirements

Handout	_____
MCQ	_____
Statistical Tables	_____
Graph Paper	_____
Log Graph Paper	_____
Other Material	_____

No. of Pages	_____
Department(s)	_____

1. (a) Write each of the following statements using either the Floyd Hoare triple $\vdash [P]C[Q]$ or $\vdash \{P\}C\{Q\}$ and the logical constants T for true and F for false:
- (i) “If the precondition holds, the code will terminate”
 - (ii) “If the code terminates, the postcondition will hold”
 - (iii) “If the precondition holds, the postcondition will hold”
 - (iv) “The code will terminate”
 - (v) “The code will terminate and the postcondition will hold”
- (as an example, “The code will not terminate” is $\vdash \{T\}C\{F\}$). (15 marks)

(b) Explain the distinction between

- (i) Specification and Implementation
- (ii) An Axiom and an Inference Rule
- (iii) An Axiom and a Theorem

(9 marks)

(c) Suppose we have proven $\vdash \{P\}C_1\{Q\}$, $\vdash [Q]C_2[R]$ and $\vdash [R]C_3[S]$. From these, which of the following can we conclude is true?

- (i) $\vdash \{P\}C_1; C_2\{R\}$
- (ii) $\vdash [P]C_1; C_2[R]$
- (iii) $\vdash \{P\}C_1; C_2; C_3\{S\}$
- (iv) $\vdash \{Q\}C_2; C_3\{S\}$
- (v) $\vdash \{S\}C_3; C_2; C_1\{P\}$
- (vi) $\vdash \{R\}C_3\{S\}$

(6 marks)

2. (a) Prove the derived inference rules

(i)

$$\frac{\vdash \{P\} C \{Q_1\} \quad \vdash \{P\} C \{Q_2\}}{\vdash \{P\} C \{Q_1 \leftrightarrow Q_2\}}$$

(ii)

$$\frac{\vdash \{P_1\} C \{Q_1\} \quad \vdash \{P_2\} C \{Q_2\}}{\vdash \{P_1 \wedge P_2\} C \{\neg Q_1 \rightarrow Q_2\}}$$

(15 marks)

(b) In each of the following state which (if either) of the two conditions is weaker.

- (i) $xy > 0$; $(x > 0) \wedge (y > 0)$
- (ii) $x > 2$; $x \neq 2$
- (iii) $A \wedge (B \vee C)$; $C \vee (A \wedge B)$
- (iv) $x < 6$; $x \neq 5$
- (v) $(\forall_x)P(x) = 7$; $(\exists_y)P(y) > 6$

(15 marks)

3. (a) Prove the following formulas of Floyd Hoare logic:

- (i) $\{(2 \times X) > 2\} X := X - 1; X := X + X \{X \geq 0\}$
- (ii) $\{True\} A := Y; Z := 1 \{Z = X^{Y-A}\}$
- (iii) $\{Y = n \wedge X = m\}$
 $X := X + Y; Y := X - Y; X := X - Y$
 $\{X = n \wedge Y = m\}$

(18 marks)

(b) For all variables in the following expressions, state whether they are free or bound:

- (i) $(\forall x. (\exists y. P(x) \Rightarrow Q(y)) \wedge \neg R(z, x))$
- (ii) $P(x, y) \vee (\exists z. Q(x)) \Rightarrow R(z)$
- (iii) $\forall x. ((y \geq 1) \wedge (x \geq y)) \wedge \forall z. (\neg(z = 1))$

(12 marks)

4. (a) Calculate the results of carrying out the following substitutions:

- (i) $(X^3 - 3 * X)[X - 1/X]$
- (ii) $(X + Y)[Y, X/X, Y]$
- (iii) $(X + Y)[X/Y][Y/X]$
- (iv) $(Z \leq Y)[X + Z/Z][X + Z, Y/Y, X]$

(12 marks)

(b) Prove:

$\vdash [Y > 0]$
 $R := X; Q := 0;$
WHILE $Y \leq R$ **DO**
 BEGIN $R := R - Y; Q := Q + 1$ **END**
 $[X = Y \times Q + R \wedge R < Y]$

(N.B.: this is a **total** correctness specification).

(18 marks)

5. (a) Prove:

$\vdash \{A[X] = x \wedge A[Y] = y \wedge X \neq Y\}$
 $A[X] := A[X] + A[Y];$
 $A[Y] := A[X] - A[Y];$
 $A[X] := A[X] - A[Y]$
 $\{A[X] = y \wedge A[Y] = x\}$

(15 marks)

(b) Using Weakest Precondition Semantics, determine

$$WP(y := 5; x := x - 1; x := -x * y, x > 10).$$

(9 marks)

(c) In Program Refinement suppose $\#[P, Q_1] = 5$, $\#[Q_1, R] = 6$, $\#[P, Q_2] = 7$, and $\#[Q_2, R] = 8$. Determine a lower bound for

$$\#\{C \mid \vdash [P]C[R]\}.$$

(6 marks)

Supplementary Material – Axioms and Inference Rules

Assignment Axiom $\vdash \{P[E/I]\} I := E \{P\}$

Array Assignment Axiom $\vdash \{P[A\{E_1 \leftarrow E_2\}/A]\} A[E_1] := E_2 \{P\}$

Precondition Strengthening $\frac{\vdash P \Rightarrow Q \quad \vdash \{Q\} C \{R\}}{\vdash \{P\} C \{R\}}$

Postcondition Weakening $\frac{\vdash \{P\} C \{Q\} \quad \vdash Q \Rightarrow R}{\vdash \{P\} C \{R\}}$

Sequencing $\frac{\vdash \{P\} C_1 \{Q\} \quad \vdash \{Q\} C_2 \{R\}}{\vdash \{P\} C_1; C_2 \{R\}}$

Blocks $\frac{\vdash \{P\} C \{Q\}}{\vdash \{P\} \text{BEGIN } VAR I_1; \dots VAR I_n; C \text{END} \{Q\}}$

Two-Armed Conditional $\frac{\vdash \{P \wedge S\} C_1 \{Q\} \quad \vdash \{P \wedge \neg S\} C_2 \{Q\}}{\vdash \{P\} \text{IF } S \text{ THEN } C_1 \text{ ELSE } C_2 \{Q\}}$

One-Armed Conditional $\frac{\vdash \{P \wedge S\} C \{Q\} \quad \vdash (P \wedge \neg S) \Rightarrow Q}{\vdash \{P\} \text{IF } S \text{ THEN } C \{Q\}}$

While $\frac{\vdash \{P \wedge S\} C \{P\}}{\vdash \{P\} \text{WHILE } S \text{ DO } C \{P \wedge \neg S\}}$

Specification Conjunction $\frac{\vdash \{P_1\} C \{Q_1\} \quad \vdash \{P_2\} C \{Q_2\}}{\vdash \{P_1 \wedge P_2\} C \{Q_1 \wedge Q_2\}}$

Specification Disjunction $\frac{\vdash \{P_1\} C \{Q_1\} \quad \vdash \{P_2\} C \{Q_2\}}{\vdash \{P_1 \vee P_2\} C \{Q_1 \vee Q_2\}}$