

Design Theory and Algebra

Warwick de Launey

IDA/CCR

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Outline

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New Book

de Launey and Flannery, Design Theory and Algebra.

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What's in this talk

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- some ideas from the book

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- some ideas from the book
- some sample results

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What's in this talk

- some ideas from the book
- some sample results
- a few words of encouragement

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Combinatorial mathematics is tremendously alive at this moment, and we believe that its greatest truths are still to be revealed.

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Important Observation

A design is a purely combinatorial concept.

Combinatorial Constraints

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A. Define the set of pairs of rows that are allowed to appear together in an array.

Orthogonality Set

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Pairwise Combinatorial Design

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A $\text{PCD}(v, \Lambda)$ is a $v \times b$ array whose rows are pairwise Λ -orthogonal.

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Proposition

An error-correcting code with minimum distance at least d and block-length n is a $\text{PCD}(\Lambda_{\text{BC}(n,d)})$.

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A Latin Square is a $\text{PCD}(\Lambda_{\text{LS}(n)})$.

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A Hadamard matrix is a $\text{PCD}(\Lambda_{\text{H}(n)})$.

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- Are there circulant $\text{PCD}(\Lambda)$?

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- $v_{\max}(\Lambda_{\text{BC}(n,d)}) = ??$

Almost-Hadamard Matrices Exist

Almost-Half-Hadamard Matrices Exist

Theorem (de Launey and Gordon)

For all sufficiently large n , $v_{\max}(\Lambda_{H(n)}) \geq \frac{1}{2}n + O(n^{113/132+o(1)})$.

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- (Graham and Shparlinski) Fairly short arithmetic sequences modulo 2^a always contain a prime power.
- So we can paste two Hadamard matrices of orders close to $n/2$ together.



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- if $H = [h_{ij}]$ is a generalized Hadamard matrix, then so is $H^* = [h_{ij}^*]$.
- if A is $(1, -1)$ -matrix with all inner products equal to ± 1 , then so is A^T .

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Conjecture (de Launey, Horadam)

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Conjecture (Ito)

There is a relative difference set in the generalized quaternion group Q_{8t} of order $8t$ for every t .

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- (with Michael Smith) $a = 8$.
- (with Hadi Kharaghani) $c = 10$, $a = 4/5$.
- for any $\epsilon > 0$, we may take $a = \epsilon$ for a positive density of g 's.

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Conjecture

There are enough complex complementary sequences with zero autocorrelation so that the WDL-Kharaghani construction gives α as small as you like.

So Where's the Algebra?



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Key Idea

Turn the arrays into matrices.

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Ambient Rings

Embed \mathcal{A} into a ring \mathcal{R} where the arithmetic in \mathcal{R} is compatible with Λ -orthogonality.

Where We Are Heading

Modeling Λ -Equivalence

\mathcal{R} contains a row group R , a column C such that two $\text{PCD}(\Lambda)$ designs A and B are Λ -equivalent if and only if there exist $P \in \text{Mon}(v, R)$ and $Q \in \text{Mon}(v, C)$ such that

$$A = P^* B Q.$$

Gram Property

A is a $\text{PCD}(\Lambda)$ if and only if $AA^* \in \text{Gram}_{\mathcal{R}}(\Lambda)$.

The Gram Property: Involutory Rings

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- 3 $(r + s)^* = r^* + s^*$.

The Gram Property II: Definition

Definition

Λ has the Gram Property over the involutory ring \mathcal{R} if any $v \times v$ $(0, \mathcal{A})$ -array A is a $\text{PCD}(\Lambda)$ if and only if for some finite set \mathcal{B} ,

$$AA^* = B \quad \text{for some } B \in \mathcal{B}.$$

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- 1 \mathcal{R} is said to have the Gram Property for Λ ,
- 2 The set \mathcal{B} is denoted $\text{Gram}_{\mathcal{R}}(\Lambda)$.

Λ -Equivalence I: Two Groups

The Groups $\Pi_{\Lambda}^{\text{row}}$ and $\Pi_{\Lambda}^{\text{col}}$

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- 1 $\Pi_{\Lambda}^{\text{row}}$ is set of permutations ρ on \mathcal{A} such that Λ is invariant under the operation where each entry $x \in \mathcal{A}$ in a row is replaced by $\rho(x)$.

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- 2 $\Pi_{\Lambda}^{\text{col}}$ is set of permutations ρ on \mathcal{A} such that Λ is invariant under the operation where each entry $x \in \mathcal{A}$ in a column is replaced by $\rho(x)$.

Λ -Equivalence II: Row and Column Operations



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(row operations)

- 1 swapping two rows,
- 2 for some $\rho \in \Pi_{\Lambda}^{\text{row}}$, replacing each entry $x \in \mathcal{A}$ in a row with $\rho(x)$.

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Definition

Two $\text{PCD}(\Lambda)$ designs are equivalent if there is a sequence of row and column operations which turns one into the other.

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Note

R and C are embeddings of $\Pi_{\Lambda}^{\text{row}}$ and $\Pi_{\Lambda}^{\text{col}}$ in \mathcal{R} .

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Theorem

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Problem

Study these rings.

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- ② Solve the matrix equation $AA^* = B$ over the ambient ring \mathcal{R} ,

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- ① Study and compute Automorphism Groups of Designs,
- ② Solve the matrix equation $AA^* = B$ over the ambient ring \mathcal{R} ,
- ③ Solve the group ring equation $\alpha\alpha^* = b$ over the group ring $\mathcal{R}[G]$.

How to Proceed

- ① Work on the hard problems,
- ② Do some whimsical math,
- ③ Share your ideas,
- ④ Understand your colleagues strengths,
- ⑤ Try to find someone, who knows something that you don't know, who is willing to work with you on your problems.

My teachers (open literature):

- 1 K.T. Arasu (character theory),
- 2 David Cargo (recursions),
- 3 Dane Flannery (computational algebra),
- 4 Dan Gordon (number theory),
- 5 Kathy Horadam (cohomology),
- 6 David Levin (probability theory),
- 7 Martin Liebeck (representation theory),
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- 9 Dick Stafford (finite group theory).

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Special thanks to my advisor: Jennifer Seberry (design theory).



